

Physics for Game Programmers:

Collision Detection

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Collision Detection

- Find all pairs of objects that are colliding now, or will collide over the next frame.
- Compute data for response:
 - Contact normal
 - Contact point
 - Penetration depth

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The Problem







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The Solution



Construct Plausible Trajectories

- Limited to trajectories involving piecewise constant linear velocities.
- Angular velocities are ignored. Rotations are considered instantaneous.

No Continuous Rotations?

- Solving continuous rotations is a lot trickier, so we dodge the issue.
- Tunneling may occur for rotating objects, but is less visible and often acceptable.
- Only doing continuous translations fixes our problems and is doable in real time.

Collision Objects

- Static environment (buildings, terrain) is typically modeled using polygon meshes.
- Moving objects (player, NPCs, vehicles, projectiles) are typically convex shapes.
- We need to detect convex-convex and convex-mesh collisions.

Convex Shapes



Polytopes





Quadric Shapes



Configuration Space

• The configuration space obstacle (CSO) of objects A and B is the set of all vectors from a point of B to a point of A.

$$A - B = \{\mathbf{a} - \mathbf{b} : \mathbf{a} \in A, \mathbf{b} \in B\}$$

Configuration Space (cont'd)

CSO is basically one object dilated by the other:



Translation

Translation of A and/or B results in a translation of A – B.



Rotation

Rotation of A and/or B changes the shape of A – B.





Configuration Space?

- Collision queries on a pair of convexes are reduced to queries on the position of the origin with respect to the CSO.
- Point queries are easier than queries on pairs of shapes.

Queries: Distance

• The distance between two objects is the distance from the origin to the CSO.

$$d(A,B) = \min\left\{ \left\| \mathbf{x} \right\| : \mathbf{x} \in A - B \right\}$$

Queries: Intersection Testing

 The objects intersect (have a common point) if the origin is contained by the CSO.

$A \cap B \neq \emptyset \Leftrightarrow \mathbf{0} \in A - B$

Queries: Penetration Depth

 The penetration-depth vector is the shortest translation that resolves a penetration, i.e., the point on the CSO's boundary closest to the origin.

$$p(A,B) = \inf\left\{ \left\| \mathbf{x} \right\| : \mathbf{x} \notin A - B \right\}$$

Queries: Shape Casting

 Finding collisions that occur over a frame for A translated over s and B over t boils down to a ray cast from the origin onto the CSO along the vector r = t − s.

$$\min\{\lambda: \lambda \mathbf{r} \in A - B, 0 \le \lambda \le 1\}$$

Ray Query on the CSO



Separating Axis



Separating Axis Theorem (SAT)

- For each pair of disjoint polytopes, of which at least one has a volume, there exists a separating axis that is orthogonal to:
- a face of either polytope, or
- an edge from each polytope

SAT Sketchy Proof

- The CSO of polytopes is a polytope and has a volume.
- For disjoint polytopes, the origin lies on the outside of at least one face of the CSO.
- A face of the CSO is either the CSO of a face and a vertex or of two edges.

- Test all face normals and all cross products of edge directions.
- If none of these vectors yield a separating axis then the polytopes must intersect.
- Given polytopes with resp. f_1 and f_2 faces and e_1 and e_2 edge directions, we need to test at most $f_1 + f_2 + e_1 * e_2$ axes.

Polytope 1 Polytope 2 #Axes

Line segment Box Triangle Box Box Box Tetrahedron Tetrahedron

0 + 3 + 1 + 3 = 6

Separating Axis Method

Polytope 1 Polytope 2 #Axes

Line segment Box Triangle Box Box Box Tetrahedron Tetrahedron

Polytope 1 Polytope 2 #Axes

Line segmentBox0+3+1*3=6TriangleBox1+3+3*3=13BoxBoxTetrahedron

Polytope 1 Polytope 2 #Axes

Line segmentBox0+3+1*3=6TriangleBox1+3+3*3=13BoxBox3+3+3*3=15TetrahedronTetrahedron

Polytope 1 Polytope 2 #Axes

Line segmentBox0 + 3 + 1*3 = 6TriangleBox1 + 3 + 3*3 = 13BoxBox3 + 3 + 3*3 = 15TetrahedronTetrahedron4 + 4 + 6*6 = 44

Separating Axis Queries

- Suitable for intersection testing, most notably in bounding box hierarchies.
- Too expensive for general polytopes due to $O(n^3)$ complexity.
- In case of intersection, the axis for which overlap is shallowest is a proper direction for the penetration depth vector.

GJK Does It All

- GJK is an iterative method that computes closest points.
- The GJK ray cast can perform continuous collision detection.
- The *expanding polytope algorithm* (EPA) returns the penetration-depth vector.

GJK Algorithm

- Approximate the point of the CSO closest to the origin by generating a sequence of simplices inside the CSO.
- A *simplex* is a point, a line segment, a triangle, or a tetrahedron.
- Each new simplex lies closer to the origin than its predecessor.

GJK Algorithm (cont'd)

- Simplex vertices are computed using support mappings. (Definition follows.)
- Terminate as soon as the current simplex is close enough.
- In case of an intersection, the simplex contains the origin.

Support Mappings

 A support mapping s_A of an object A maps a vector v to a point of A that lies furthest in the direction of v.

$$\mathbf{v} \cdot s_A(\mathbf{v}) = \max\left\{\mathbf{v} \cdot \mathbf{x} : \mathbf{x} \in A\right\}$$

Support Mappings


Affine Transformation

 Shapes can be translated, rotated, and scaled. For T(x) = Bx + c, we have

$$s_{\mathbf{T}(A)}(\mathbf{v}) = \mathbf{T}(s_A(\mathbf{B}^{\mathrm{T}}\mathbf{v}))$$



Convex Hull

Convex hulls of arbitrary convex shapes are readily available.

$$S_{\text{conv}\{X_0,...,X_{n-1}\}}(\mathbf{v}) = S_{\{s_{X_0}(\mathbf{v}),...,s_{X_{n-1}}(\mathbf{v})\}}(\mathbf{v})$$



Minkowski Sum

• Shapes can be fattened by *Minkowski addition*.

$$s_{A+B}(\mathbf{v}) = s_A(\mathbf{v}) + s_B(\mathbf{v})$$
$$s_{A-B}(\mathbf{v}) = s_A(\mathbf{v}) - s_B(-\mathbf{v})$$
$$+ =$$

GJK Steps (1/6)

 Suppose we have a simplex inside the CSO...



GJK Steps (2/6)

 ...and the point v of the simplex closest to the origin.



GJK Steps (3/6)

• Compute support point $\mathbf{w} = s_{A-B}(-\mathbf{v})$.



GJK Steps (4/6)

Add support point w to the current simplex.



GJK Steps (5/6)

Compute the closest point v' of the new simplex.



GJK Steps (6/6)

 Discard all vertices that do not contribute to v'.



Separating Axis

- If only an intersection test is needed then let GJK terminate as soon as the lower bound v·w becomes positive.
- For a positive lower bound v·w, the vector v is a separating axis.

Separating Axis (cont'd)

• The supporting plane through w separates the origin from the CSO.



Separating Axes and Coherence

- Separating axes can be cached and reused as initial v in future tests on the same object pair.
- When the degree of frame coherence is high, the cached v is likely to be a separating axis in the new frame as well.
- An incremental version of GJK takes roughly one iteration per frame for smoothly moving objects.

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GJK Ray Cast



GJK Ray Cast

- Do a standard GJK iteration, and use the support planes as clipping planes.
- Each time the ray is clipped, the clip point λr becomes the new origin.
- ...and the new simplex is the last-found support point w wrt the new origin.
- The normal -v of the last clipping plane is the normal at the hit point.

GJK Ray Cast



- Accuracy can be traded for performance by tweaking the error tolerance ε_{tol} .
- A greater tolerance results in fewer iterations but less accurate hit points and normals.

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Accuracy vs. Performance

 $\varepsilon_{tol} = 10^{-7}$, avg. time: 3.65 µs @ 2.6 GHz



 $\varepsilon_{tol} = 10^{-6}$, avg. time: 2.80 µs @ 2.6 GHz



 $\varepsilon_{tol} = 10^{-5}$, avg. time: 2.03 µs @ 2.6 GHz



 $\varepsilon_{tol} = 10^{-4}$, avg. time: 1.43 µs @ 2.6 GHz



 $\varepsilon_{tol} = 10^{-3}$, avg. time: 1.02 µs @ 2.6 GHz



 $\varepsilon_{tol} = 10^{-2}$, avg. time: 0.77 µs @ 2.6 GHz



$\varepsilon_{tol} = 10^{-1}$, avg. time: 0.62 µs @ 2.6 GHz



GJK Algorithm: Pros

- Extremely versatile:
 - Applicable to any combination of convex shape types.
 - Computes distances, common points, and separating axes.
 - Can be tailored for finding space-time collisions.
 - Allows a smooth trade-off between accuracy and speed.

GJK Algorithm: Pros (cont'd)

- Performs well:
 - Exploits frame coherence.
 - Competitive with dedicated solutions for polytopes (Lin-Canny, V-Clip, SWIFT) .
- Despite its conceptual complexity, implementing GJK is not too difficult.
- Small code size.

GJK Algorithm: Cons

- Difficult to grasp:
 - Concepts from linear algebra and convex analysis (determinants, Minkowski addition), take some time to get comfortable with.
 - Maintaining a "geometric" mental image of the workings of the algorithm is challenging and not very helpful.

GJK Algorithm: Cons (cont'd)

- Suffers from numerical issues:
 - Termination is governed by predicates that rely on tolerances.
 - Despite the use of tolerances, certain "hacks" are needed in order to guarantee termination in all cases.
 - Using 32-bit floating-point numbers is doable but tricky.

Resting Contacts

- Contact data for resting contacts are obtained through a hybrid approach.
- Objects are dilated slightly to add a skin.
- For interpenetrations that are only skindeep the closest points of the "bones" give us the contact data.

Shallow Interpenetrations



Resting Contacts

- For deeper interpenetrations contact data are obtained from the penetrationdepth vector.
- This should only be necessary in emergencies.

Deep Interpenetrations



Meshes Have Bumpy Edges



Solving Bumpy Edges

- Obtain *barycentric coordinates* of the closest point returned by GJK.
- Use these coordinates to interpolate the vertex normals.
- Similar to Phong shading: Use a normalized lerp.

Smooth Interpolated Normals



References

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